

# Proof Nets

A graph syntax for proofs

MLL

$$\frac{}{\vdash A^\perp, A} ax \quad \frac{\vdash \Gamma, A \quad \vdash \Delta, A^\perp}{\vdash \Gamma, \Delta} cut$$

$$\frac{\vdash \Gamma, A \quad \vdash \Delta, B}{\vdash \Gamma, \Delta, A \otimes B} \otimes \quad \frac{\vdash \Gamma, A, B}{\vdash \Gamma, A \wp B} \wp$$

Reference:  
Notes on proof-nets by Olivier Laurent

(Note: most slides are taken from the notes of Olivier Laurent)

Recall that linear negation is defined :

$$(X^\perp)^\perp = X$$

$$(A \otimes B)^\perp = A^\perp \wp B^\perp$$

$$(A \wp B)^\perp = A^\perp \otimes B^\perp$$

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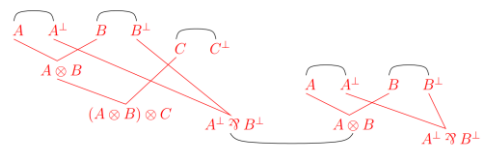
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## Forgetting Sequential Structure

$$\frac{\frac{\frac{}{\vdash A, A^\perp} (ax) \quad \frac{}{\vdash B, B^\perp} (ax)}{\vdash A \otimes B, A^\perp, B^\perp} (\otimes) \quad \frac{}{\vdash C, C^\perp} (ax)}{\vdash (A \otimes B) \otimes C, A^\perp, B^\perp, C^\perp} (\otimes) \quad \frac{}{\vdash A \otimes B, A^\perp, B^\perp} (ax)}{\vdash (A \otimes B) \otimes C, A^\perp \wp B^\perp, C^\perp} (\wp) \quad \frac{}{\vdash A \otimes B, A^\perp, B^\perp} (ax)}{\vdash (A \otimes B) \otimes C, A^\perp, B^\perp, C^\perp} (cut) \quad \frac{}{\vdash (A \otimes B) \otimes C, A^\perp, B^\perp, C^\perp} (\wp)}{\vdash (A \otimes B) \otimes C, A^\perp \wp B^\perp, C^\perp} (\wp)$$

## Forgetting Sequential Structure

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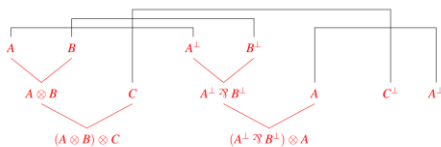
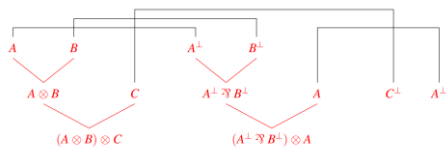
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## Forgetting Sequential Structure

$$\begin{array}{l} (1) \frac{}{\vdash A, A^\perp} ax \quad (2) \frac{}{\vdash B, B^\perp} ax \\ (3) \frac{}{\vdash A \otimes B, A^\perp, B^\perp} \otimes \quad (4) \frac{}{\vdash C, C^\perp} ax \\ (5) \frac{}{\vdash (A \otimes B) \otimes C, A^\perp, B^\perp, C^\perp} \otimes \\ (6) \frac{}{\vdash (A \otimes B) \otimes C, A^\perp \wp B^\perp, C^\perp} \wp \quad (7) \frac{}{\vdash A, A^\perp} ax \\ (8) \frac{}{\vdash (A \otimes B) \otimes C, (A^\perp \wp B^\perp) \otimes A, C^\perp, A^\perp} \otimes \end{array}$$

## Forgetting Sequential Structure

$$\begin{array}{l} (1) \frac{}{\vdash A, A^\perp} ax \quad (2) \frac{}{\vdash B, B^\perp} ax \\ (3) \frac{}{\vdash A \otimes B, A^\perp, B^\perp} \otimes \quad (4) \frac{}{\vdash C, C^\perp} ax \\ (5) \frac{}{\vdash (A \otimes B) \otimes C, A^\perp, B^\perp, C^\perp} \otimes \\ (6) \frac{}{\vdash (A \otimes B) \otimes C, A^\perp \wp B^\perp, C^\perp} \wp \quad (7) \frac{}{\vdash A, A^\perp} ax \\ (8) \frac{}{\vdash (A \otimes B) \otimes C, (A^\perp \wp B^\perp) \otimes A, C^\perp, A^\perp} \otimes \end{array}$$



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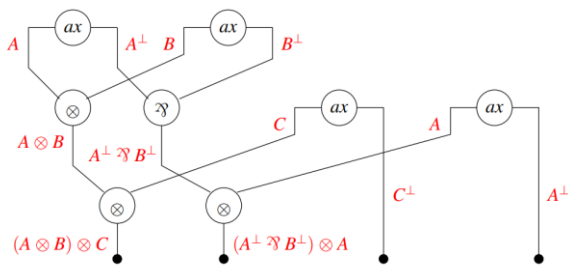
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# Proof Nets

A graph syntax for proofs

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In the graphical representation of a proof structure, we do not mention explicitly the direction of edges, but we draw them in such a way that direction is represented in a top-down way:

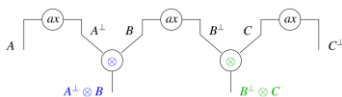


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## example

Translate each of these sequent calculus proofs. Start from axioms....

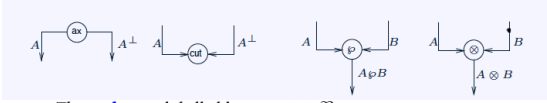
$$\frac{\frac{\frac{}{\vdash A, A^\perp} ax}{} \quad \frac{\frac{}{\vdash B, B^\perp} ax}{} \quad \frac{}{\vdash C, C^\perp} ax}{\vdash A, A^\perp \otimes B, B^\perp \otimes C, C^\perp} \otimes \quad \frac{\frac{}{\vdash A, A^\perp} ax}{} \quad \frac{\frac{\frac{}{\vdash B, B^\perp} ax}{} \quad \frac{}{\vdash C, C^\perp} ax}}{\vdash B, B^\perp \otimes C, C^\perp} \otimes}{\vdash A, A^\perp \otimes B, B^\perp \otimes C, C^\perp} \otimes$$



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## Proof structures

A **proof structure**  $M$  is a **labelled directed acyclic graph (DAG)** with possibly pending edges (i.e. some edges may have no source and/or no target) built over the alphabet of nodes which is represented below.  
(Note: in figures, the edges orientation is always top-bottom.)



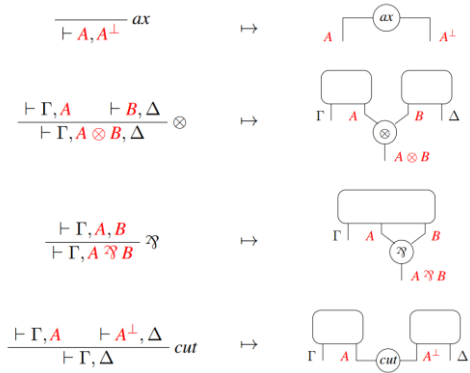
- The **nodes** are labelled by ax, cut,  $\otimes$ ,  $\cap$
- The **edges** are labelled by MLL formulas.

For each node/link: **premises** = entering edges, **conclusions** = exiting edges

The **conclusions** of  $M$  is the set of pending edges of  $M$ .

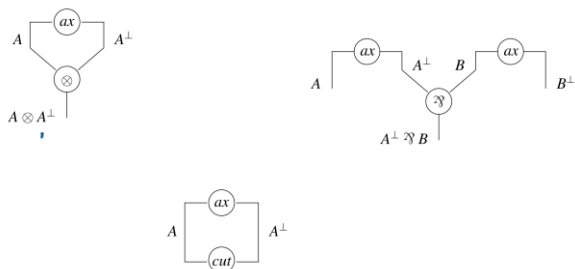
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## From proofs to proof structures



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## Is every structure the image of an MLL proof?



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Proof Nets

**A PROOF NET** is a proof structure which is the image of an MLL proof

*Internal condition!*

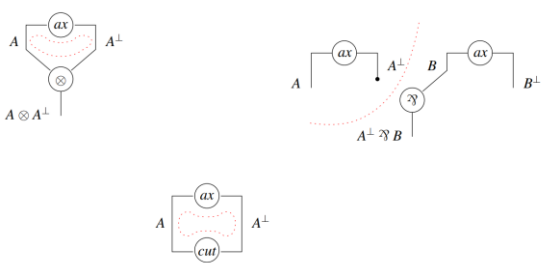
Purely geometrical conditions (correction) characterize the proof structures which are proof nets

Theorem 1. A proof structure is correct iff it is a proof net.

Correctness criteria:

- (LT) Long trip [Girard]
- (AC) Connected-Acyclic [Danos-Regnier]
- ...

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**Acyclicity.** A multiplicative proof structure is *acyclic* if its switching graphs do not contain any undirected cycle.

A proof structure with  $p$   $\bowtie$  nodes induces  $2^p$  switchings and thus  $2^p$  switching graphs. A switching graph is not a proof structure in general since its  $\bowtie$  nodes have only one premise.

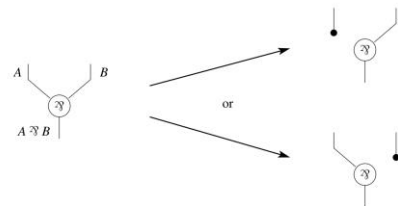
A *connected component* of a switching graph is a connected component of its underlying (undirected) multigraph.

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Danos-Regnier Criterion

Correctness Criterion

Switching Graphs



Correctness

- Switching graphs are **acyclic**.
- Switching graphs are **connected**.

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Definition 2 (Correctness criterion AC (Danos-Regnier)). Let  $R$  be a proof structure.

A *switching*  $s$  is a function on the nodes of  $R$ , which chooses, for each  $\bowtie$ -link, either the left or the right premise.

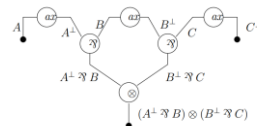
A proof structure  $R$  is *correct* if for each switching, the unoriented graph obtained by erasing for each  $\bowtie$ -link of  $R$  the edges not chosen by  $s$ , is:

*connected and acyclic*

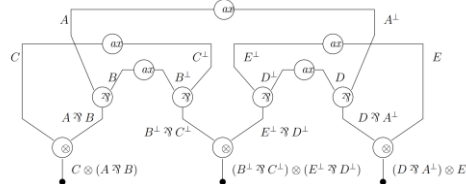
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Is this correct?

PN1:



PN2:



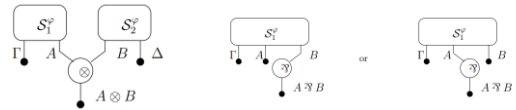
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- Correctness guarantees:
- ✓ Graph is image of a proof (sequentialization)
- ✓ Normalization progresses (no deadlocks)
- ✓ Normalization terminates (no infinite cycles)

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**Soundness**

**Proposition 4.1.1** (Soundness of Correctness). *The translation of a sequent calculus proof of MLL is a connected multiplicative proof net.*



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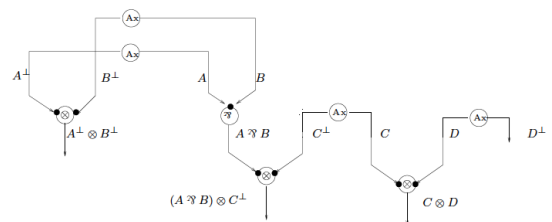
Sequentialization

**Theorem 4.1.1** (Sequentialization). *Any connected multiplicative proof net is the translation of a sequent calculus proof of MLL.*

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Sequentialization answers the question:

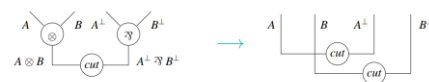
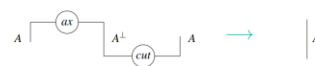
We have a proof net. The problem: it is the image of a sequent calculus proof? And which?



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Normalization  
(local graph reductions!)

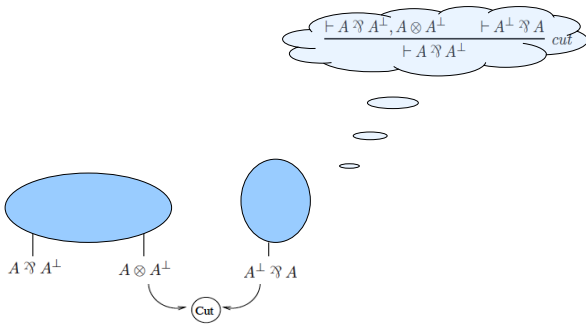
The beauty of proof nets is normalization



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Let us try out!



Write a proof net with this conclusion... and normalize it

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Let us try one more. First, write a proof net with this conclusion...

$$(X \otimes X) \multimap (X \otimes X) = (X \otimes X)^\perp \multimap (X \otimes X) = (X^\perp \multimap X^\perp) \multimap (X \otimes X)$$

TIP: How we write a proof net? As before, all proof nets with the same conclusion, start with the same nodes (the formula tree!) What distinguishes different proofs are the axiom links

To distinguish the different occurrences of atoms, let us write indices:

$$(X_1^\perp \multimap X_2^\perp) \multimap (X_3 \otimes X_4)$$

In this case, we have two possible proofs, corresponding to two possible way to write axioms:

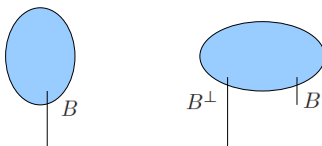
- 1,3 and 2,4
- OR
- 1,4 and 2,3

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When we have a formula whose normal proofs are exactly two, we have a good candidate to code BOOLEANS :)

Let us indicate the formula  $(X^\perp \multimap X^\perp) \multimap (X \otimes X)$  with B (for boolean). We call one proof **true**, and the other **false**...

We can feed one of our two values to a proof which takes a boolean, and return a boolean.



We know that the normal form (i.e the result of computation) will be of type B... Hence one of our two values.

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How we write a proof net of these conclusions?

$A \multimap A^\perp$  must type an edge conclusion of a par link, with premisses ....

$A \otimes A^\perp$  must type an edge conclusion of a tensor link, with premisses ....

Then we have to choose the axiom links!

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In sequent calculus, they correspond to these two proofs (one uses exchange, one no)

$$\frac{\frac{\frac{\frac{\frac{\vdash X_3^\perp, X_1}{\vdash X_1^\perp, X_2^\perp, X_3 \otimes X_4}}{\vdash X_1^\perp \multimap X_2^\perp, X_3 \otimes X_4}}{\vdash X_1^\perp \multimap X_2^\perp, X_3 \otimes X_4}}{\vdash (X_1 \otimes X_3)^\perp \multimap (X_2 \otimes X_4)}}{\vdash X_3^\perp, X_2}{\vdash X_4^\perp, X_1} \otimes \quad \frac{\frac{\frac{\frac{\frac{\vdash X_3^\perp, X_2}{\vdash X_1^\perp, X_2^\perp, X_3 \otimes X_4}}{\vdash X_1^\perp \multimap X_2^\perp, X_3 \otimes X_4}}{\vdash X_1^\perp \multimap X_2^\perp, X_3 \otimes X_4}}{\vdash (X_1 \otimes X_3)^\perp \multimap (X_2 \otimes X_4)}}{\vdash X_4^\perp, X_1} \otimes$$

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Try to normalize one of the proofs of  $(X_1^\perp \multimap X_2^\perp) \multimap (X_3 \otimes X_4)$

with the proof net which has conclusions

$$(X_1 \otimes X_2) \otimes (X_3^\perp \multimap X_4^\perp) \quad (X_5^\perp \multimap X_6^\perp) \multimap (X_7 \otimes X_8)$$

and axiom links: (1,6) (2,5) (3,7) (4,8)

What is the function coded by this proof net?

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