

# Linear Logic

Reference:  
Lecture Notes by Olivier Laurent (in French)

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## Recall LK ?

**Groupe logique multiplicatif.**

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma' \vdash B, \Delta'}{\Gamma, \Gamma' \vdash A \wedge B, \Delta, \Delta'} \wedge^{\text{mul}}R \quad \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \wedge^{\text{mul}}L$$

$$\frac{\Gamma \vdash A, B, \Delta}{\Gamma \vdash A \vee B, \Delta} \vee^{\text{mul}}R \quad \frac{\Gamma, A \vdash \Delta \quad \Gamma', B \vdash \Delta'}{\Gamma, \Gamma', A \vee B \vdash \Delta, \Delta'} \vee^{\text{mul}}L$$

**Groupe logique additif.**

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma' \vdash B, \Delta}{\Gamma \vdash A \wedge B, \Delta} \wedge^{\text{add}}R \quad \frac{\Gamma, A \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \wedge^{\text{add}}L_1 \quad \frac{\Gamma, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \wedge^{\text{add}}L_2$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma \vdash A \vee B, \Delta} \vee^{\text{add}}R_1 \quad \frac{\Gamma \vdash B, \Delta}{\Gamma \vdash A \vee B, \Delta} \vee^{\text{add}}R_2 \quad \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \vee B \vdash \Delta} \vee^{\text{add}}L$$

**Groupe structurel.**

$$\frac{\Gamma \vdash \Delta}{\sigma(\Gamma) \vdash \tau(\Delta)} \text{ctr}$$

$$\frac{\Gamma, A, A \vdash \Delta}{\Gamma, A \vdash \Delta} \text{crl} \quad \frac{\Gamma \vdash A, A, \Delta}{\Gamma \vdash A, \Delta} \text{crr}$$

$$\frac{\Gamma \vdash \Delta}{\Gamma, A \vdash \Delta} \text{wll} \quad \frac{\Gamma \vdash \Delta}{\Gamma \vdash A, \Delta} \text{wrr}$$

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# LL Sequent Calculus

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**Formulas:**

$$A ::= X \mid A \wp A \mid A \& A \mid ?A \mid X^\perp \mid A \otimes A \mid A \oplus A \mid !A$$

**Groupe multiplicatif.**

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma' \vdash B, \Delta'}{\Gamma, \Gamma' \vdash A \otimes B, \Delta, \Delta'} \otimes R \quad \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \otimes B \vdash \Delta} \otimes L$$

$$\frac{\Gamma \vdash A, B, \Delta}{\Gamma \vdash A \wp B, \Delta} \wp R \quad \frac{\Gamma, A \vdash \Delta \quad \Gamma', B \vdash \Delta'}{\Gamma, \Gamma', A \wp B \vdash \Delta, \Delta'} \wp L$$

**Groupe additif.**

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \& B, \Delta} \& R \quad \frac{\Gamma, A \vdash \Delta}{\Gamma, A \& B \vdash \Delta} \& L_1 \quad \frac{\Gamma, B \vdash \Delta}{\Gamma, A \& B \vdash \Delta} \& L_2$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma \vdash A \oplus B, \Delta} \oplus R_1 \quad \frac{\Gamma \vdash B, \Delta}{\Gamma \vdash A \oplus B, \Delta} \oplus R_2 \quad \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \oplus B \vdash \Delta} \oplus L$$

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**Groupe identité.**

$$\frac{}{A \vdash A} \text{ax} \quad \frac{\Gamma \vdash A, \Delta \quad \Gamma', A \vdash \Delta'}{\Gamma, \Gamma' \vdash \Delta, \Delta'} \text{cut}$$

**Groupe négation.**

$$\frac{\Gamma, A \vdash \Delta}{\Gamma \vdash A^\perp, \Delta} (\cdot)^\perp R \quad \frac{\Gamma \vdash A, \Delta}{\Gamma, A^\perp \vdash \Delta} (\cdot)^\perp L$$

**Groupe exponentiel.**

$$\frac{\Gamma \vdash A, \Delta}{\Gamma \vdash ?A, \Delta} ?R \quad \frac{\Gamma, A \vdash ?\Delta}{\Gamma, ?A \vdash ?\Delta} ?L$$

$$\frac{\Gamma \vdash A, ?\Delta}{\Gamma \vdash !A, ?\Delta} !R \quad \frac{\Gamma, A \vdash \Delta}{\Gamma, !A \vdash \Delta} !L$$

$$\frac{\Gamma \vdash ?A, ?A, \Delta}{\Gamma \vdash ?A, \Delta} \text{ctrR} \quad \frac{\Gamma, !A, !A \vdash \Delta}{\Gamma, !A \vdash \Delta} \text{ctrL}$$

$$\frac{\Gamma \vdash \Delta}{\Gamma \vdash ?A, \Delta} \text{wrr} \quad \frac{\Gamma \vdash \Delta}{\Gamma, !A \vdash \Delta} \text{wll}$$

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## MALL (no structural rules)

**Groupe multiplicatif.**

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma' \vdash B, \Delta'}{\Gamma, \Gamma' \vdash A \otimes B, \Delta, \Delta'} \otimes R \quad \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \otimes B \vdash \Delta} \otimes L$$

$$\frac{\Gamma \vdash A, B, \Delta}{\Gamma \vdash A \wp B, \Delta} \wp R \quad \frac{\Gamma, A \vdash \Delta \quad \Gamma', B \vdash \Delta'}{\Gamma, \Gamma', A \wp B \vdash \Delta, \Delta'} \wp L$$

**Groupe additif.**

$$\frac{\Gamma \vdash A, \Delta \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \& B, \Delta} \& R \quad \frac{\Gamma, A \vdash \Delta}{\Gamma, A \& B \vdash \Delta} \& L_1 \quad \frac{\Gamma, B \vdash \Delta}{\Gamma, A \& B \vdash \Delta} \& L_2$$

$$\frac{\Gamma \vdash A, \Delta}{\Gamma \vdash A \oplus B, \Delta} \oplus R_1 \quad \frac{\Gamma \vdash B, \Delta}{\Gamma \vdash A \oplus B, \Delta} \oplus R_2 \quad \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \oplus B \vdash \Delta} \oplus L$$

**Groupe négation.**

$$\frac{\Gamma, A \vdash \Delta}{\Gamma \vdash A^\perp, \Delta} (\cdot)^\perp R \quad \frac{\Gamma \vdash A, \Delta}{\Gamma, A^\perp \vdash \Delta} (\cdot)^\perp L$$

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Montrer que l'on a les distributivités suivantes :

$$A \otimes (B \oplus C) \dashv\vdash (A \otimes B) \oplus (A \otimes C)$$

$$A \wp (B \& C) \dashv\vdash (A \wp B) \& (A \wp C)$$

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Define negation:

$$A ::= X \mid A \wp A \mid A \& A$$

$$\mid X^\perp \mid A \otimes A \mid A \oplus A$$

$$(X^\perp)^\perp = X$$

$$(A \otimes B)^\perp = A^\perp \wp B^\perp$$

$$(A \wp B)^\perp = A^\perp \otimes B^\perp$$

$$(A \& B)^\perp = A^\perp \oplus B^\perp$$

$$(A \oplus B)^\perp = A^\perp \& B^\perp$$

ce qui donne  $A^{\perp\perp} = A$ .

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et on utilise des séquents de la forme  $\vdash \Gamma$ .

$$\frac{}{\vdash A^\perp, A} ax \quad \frac{\vdash \Gamma, A \quad \vdash \Delta, A^\perp}{\vdash \Gamma, \Delta} cut$$

$$\frac{\vdash \Gamma, A \quad \vdash \Delta, B}{\vdash \Gamma, \Delta, A \otimes B} \otimes \quad \frac{\vdash \Gamma, A, B \quad \vdash \Delta, A \wp B}{\vdash \Gamma, A \wp B} \wp$$

$$\frac{\vdash \Gamma, A \quad \vdash \Gamma, B}{\vdash \Gamma, A \& B} \& \quad \frac{\vdash \Gamma, A}{\vdash \Gamma, A \oplus B} \oplus_1 \quad \frac{\vdash \Gamma, B}{\vdash \Gamma, A \oplus B} \oplus_2$$

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EX  
The axiom can be restricted to atomic formulas:  
the axiom rule is admissible for every formula

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$$A \multimap B = A^\perp \wp B$$

On peut alors dériver les règles de calcul des séquents :

$$\frac{\Gamma, A \vdash B, \Delta}{\Gamma \vdash A \multimap B, \Delta} \quad \frac{\Gamma \vdash A, \Delta \quad \Gamma', B \vdash \Delta'}{\Gamma, \Gamma', A \multimap B \vdash \Delta, \Delta'}$$

$$\frac{\vdash \Gamma, A^\perp, B}{\vdash \Gamma, A \multimap B} \quad \frac{\vdash \Gamma, A \quad \vdash \Delta, B^\perp}{\vdash \Gamma, \Delta, (A \multimap B)^\perp}$$

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Élimination des coupures

On va se contenter de donner les cas clés :

$$\frac{\vdash \Gamma, A \quad \vdash A^\perp, A}{\vdash \Gamma, A} cut \rightsquigarrow \vdash \Gamma, A$$

$$\frac{\vdash \Gamma, A \quad \vdash \Delta, B \quad \vdash \Xi, A^\perp, B^\perp}{\vdash \Gamma, \Delta, A \otimes B \quad \vdash \Xi, A^\perp \wp B^\perp} \otimes \rightsquigarrow \frac{\vdash \Gamma, A \quad \vdash \Delta, B \quad \vdash \Xi, A^\perp, B^\perp}{\vdash \Gamma, \Delta, \Xi} cut \rightsquigarrow \frac{\vdash \Gamma, A \quad \vdash \Delta, \Xi, A^\perp}{\vdash \Gamma, \Delta, \Xi} cut$$

$$\frac{\vdash \Gamma, A \quad \vdash \Gamma, B \quad \vdash \Delta, A^\perp}{\vdash \Gamma, A \& B \quad \vdash \Delta, A^\perp \oplus B^\perp} \& \rightsquigarrow \frac{\vdash \Gamma, A \quad \vdash \Delta, A^\perp}{\vdash \Gamma, \Delta} cut$$

$$\frac{\vdash \Gamma, A \quad \vdash \Gamma, B \quad \vdash \Delta, B^\perp}{\vdash \Gamma, A \& B \quad \vdash \Delta, A^\perp \oplus B^\perp} \& \rightsquigarrow \frac{\vdash \Gamma, B \quad \vdash \Delta, B^\perp}{\vdash \Gamma, \Delta} cut$$

EX. 1  
Define commutative steps (just give two examples).

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### Homework Ex 1

EX. 1  
How do we define commutative steps?  
Give the following two examples:

$$\frac{\vdash \Gamma, A \quad \vdash \Delta, B}{\vdash \Gamma, \Delta, A \otimes B} \otimes$$

Cut with a formula inside  $\Gamma$ .  
How do we reduce this cut?

$$\frac{\vdash \Gamma, A \quad \vdash \Gamma, B}{\vdash \Gamma, A \& B} \&$$

Cut with a formula inside  $\Gamma$ .  
How do we reduce this cut?

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### Re-introducing the Structural Rules

In a controlled way

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$A ::= X \mid A \wp A \mid A \& A \mid \boxed{?A} \mid X^\perp \mid A \otimes A \mid A \oplus A \mid !A$

$\frac{\vdash \Gamma, A}{\vdash \Gamma, ?A} ?d$

$\frac{\vdash \Gamma, ?A, ?A}{\vdash \Gamma, ?A} ?c$

$\frac{\vdash \Gamma}{\vdash \Gamma, ?A} ?w$

$\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} !$

La règle  $?d$  ("dérélction") permet à toute formule de devenir sujette aux règles structurelles. Les règles  $?w$  ("affaiblissement") et  $?c$  ("contraction") sont les règles structurelles habituelles mais ne s'appliquent qu'aux formules dont le connecteur principal est  $?$ . La règle  $!$  ("promotion")

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### La règle promotion...

est la plus subtile. Elle permet de rendre une formule (et surtout la preuve correspondante) duplicable (ou effaçable), mais ceci nécessite un contexte adapté. On peut comprendre  $!A$  comme "A autant de fois que l'on veut". La règle écrite sous la forme :

$$\frac{\Gamma \vdash A}{\Gamma \vdash !A} !$$

dit que si  $A$  est obtenue à partir d'hypothèses utilisables autant de fois que l'on veut alors  $A$  elle-même peut être utilisée autant de fois que l'on veut. Une autre manière de comprendre la contrainte de contexte de cette règle est de regarder l'élimination des coupures.

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### Cut-elimination steps

$$\frac{\frac{\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} ! \quad \frac{\vdash \Delta, A^\perp}{\vdash \Delta, ?A^\perp} ?d}{\vdash ?\Gamma, \Delta} cut}{\vdash ?\Gamma, \Delta} cut \rightsquigarrow \frac{\vdash ?\Gamma, A \quad \vdash \Delta, A^\perp}{\vdash ?\Gamma, \Delta} cut$$

$$\frac{\frac{\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} ! \quad \frac{\vdash \Delta, ?A^\perp, ?A^\perp}{\vdash \Delta, ?A^\perp} ?c}{\vdash ?\Gamma, \Delta} cut}{\vdash ?\Gamma, \Delta} cut \rightsquigarrow \frac{\frac{\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} ! \quad \frac{\vdash \Delta, ?A^\perp, ?A^\perp}{\vdash \Delta, ?A^\perp} ?c}{\vdash ?\Gamma, \Delta} cut}{\vdash ?\Gamma, \Delta} ?c$$

$$\frac{\frac{\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} ! \quad \frac{\vdash \Delta}{\vdash \Delta, ?A^\perp} ?w}{\vdash ?\Gamma, \Delta} cut}{\vdash ?\Gamma, \Delta} cut \rightsquigarrow \frac{\vdash \Delta}{\vdash ?\Gamma, \Delta} ?w$$

$$\frac{\frac{\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} ! \quad \frac{\vdash ?\Delta, ?A^\perp, B}{\vdash ?\Delta, ?A^\perp, !B} !}{\vdash ?\Gamma, ?\Delta, !B} cut}{\vdash ?\Gamma, ?\Delta, !B} cut \rightsquigarrow \frac{\frac{\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} ! \quad \frac{\vdash ?\Delta, ?A^\perp, B}{\vdash ?\Gamma, ?\Delta, B} cut}{\vdash ?\Gamma, ?\Delta, !B} cut}{\vdash ?\Gamma, ?\Delta, !B} !$$

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### Homework Ex 2

Two formulas  $A$  and  $B$  are (linearly) equivalent, written  $A \dashv\vdash B$ , if both implications  $A \multimap B$  and  $B \multimap A$  are provable. Equivalently,  $A \dashv\vdash B$  if both  $A \vdash B$  and  $B \vdash A$  are provable.

**Ex. 2 Please pick and prove one equivalence in each of the 3 groups**

1.  $A \otimes (B \oplus C) \dashv\vdash (A \otimes B) \oplus (A \otimes C)$   
 $A \wp (B \& C) \dashv\vdash (A \wp B) \& (A \wp C)$
2.  $!(A \& B) \dashv\vdash !A \otimes !B$   
 $?(A \oplus B) \dashv\vdash ?A \wp ?B$
3.  $!A \otimes !A \dashv\vdash !A$   
 $?A \wp ?A \dashv\vdash ?A$

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## Neutral Elements

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On définit quatre éléments neutres pour les quatre connecteurs multiplicatifs et additifs : 1 (“un”) pour  $\otimes$ ,  $\perp$  (“bottom”) pour  $\wp$ ,  $\top$  (“top”) pour  $\&$  et 0 (“zéro”) pour  $\oplus$  :

$$A ::= X \mid A \wp A \mid A \& A \mid \perp \mid \top \mid ?A$$

$$\mid X^\perp \mid A \otimes A \mid A \oplus A \mid 1 \mid 0 \mid !A$$

Les règles sont obtenues comme les cas 0-aires des règles de connecteurs binaires correspondants (en particulier deux règles pour  $\oplus$  donc aucune pour 0) :

$$\frac{}{\vdash 1} \quad \frac{\vdash \Gamma}{\vdash \Gamma, \perp} \quad \frac{}{\vdash \Gamma, \top}$$

Les étapes clés d'élimination des coupures sont :

$$\frac{\frac{}{\vdash 1} \quad \frac{\vdash \Gamma}{\vdash \Gamma, \perp}}{\vdash \Gamma} \text{cut} \quad \rightsquigarrow \quad \vdash \Gamma$$

puisqu'il n'y a pas de règle pour 0.

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## On the Proof Search

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### Consequences of Cut-elimination

**Theorem 3.4.1** (cut elimination). *For every sequent  $\Gamma \vdash \Delta$ , there is a proof of  $\Gamma \vdash \Delta$  if and only if there is a proof of  $\Gamma \vdash \Delta$  that does not use the cut rule.*

**Theorem 3.4.3** (subformula property). *A sequent  $\Gamma \vdash \Delta$  is provable if and only if it is the conclusion of a proof in which each intermediate conclusion is made of subformulas of the formulas of  $\Gamma$  and  $\Delta$ .*

**Theorem 3.4.4** (consistency). *The empty sequent  $\vdash$  is not provable. Subsequently, it is impossible to prove both a formula  $A$  and its negation  $A^\perp$ ; it is impossible to prove 0 or  $\perp$ .*

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### Reversing the Reversible Rules

**Définition 22** (Connecteur réversible)  
Un connecteur  $\circ$  est réversible si lorsque  $\vdash \Gamma, A$  est dérivable avec  $\circ$  connecteur principal de  $A$ , alors  $\vdash \Gamma, A$  est dérivable avec comme dernière règle une règle d'introduction de  $\circ$ .

$$\frac{\vdash A \wp B, \Gamma}{\vdash A, B, \Gamma} (\wp^{rev}) \quad \frac{\vdash \perp, \Gamma}{\vdash \Gamma} (\perp^{rev}) \quad \frac{\vdash A_1 \& A_2, \Gamma}{\vdash A_i, \Gamma} (\&^{rev})$$

- $\frac{\frac{}{\vdash A, A^\perp} (ax) \quad \frac{}{\vdash B, B^\perp} (ax)}{\vdash A, B, A^\perp \otimes B^\perp} (\otimes) \quad \vdash A \wp B, \Gamma}{\vdash A, B, \Gamma} (cut)$
- $\frac{\frac{}{\vdash A, A^\perp} (ax)}{\vdash A, A^\perp \oplus B^\perp} (\oplus_1) \quad \vdash A \& B, \Gamma}{\vdash A, \Gamma} (cut)$

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### Reversible Rules

$$\vdash \Gamma, A \wp B \quad \rightsquigarrow \quad \frac{\frac{\frac{}{\vdash A^\perp, A} (ax) \quad \frac{}{\vdash B^\perp, B} (ax)}{\vdash A^\perp \otimes B^\perp, A, B} (\otimes) \quad \vdash \Gamma, A \wp B}{\vdash \Gamma, A, B} \wp}{\vdash \Gamma, A \wp B} \wp$$

$$\vdash \Gamma, A \& B \quad \rightsquigarrow \quad \frac{\frac{\frac{}{\vdash A^\perp, A} (ax)}{\vdash A^\perp \oplus B^\perp, A} (\oplus_1) \quad \frac{\frac{}{\vdash B^\perp, B} (ax)}{\vdash A^\perp \oplus B^\perp, B} (\oplus_2)}{\vdash \Gamma, A \& B} \& \quad \vdash \Gamma, A \& B}{\vdash \Gamma, A \& B} \&}$$

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### Reversible Rules

$$\frac{\vdash \Gamma, \perp}{\vdash \Gamma, \top} \rightsquigarrow \frac{\frac{\frac{\vdash \Gamma, \perp}{\vdash \Gamma, \perp} \text{cut}}{\vdash \Gamma, \perp} \text{cut}}{\vdash \Gamma, \top} \text{cut}$$

$$\frac{\vdash \Gamma, \top}{\vdash \Gamma, \perp} \rightsquigarrow \frac{\vdash \Gamma, \top}{\vdash \Gamma, \top} \text{cut}$$

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### Positive/negative connective

	Positive	Negative	Class
$\alpha$	atom	$A^\perp$	negation
$A \otimes B$	tensor	$A \wp B$	multiplicatives
$\mathbf{1}$	one	$\perp$	multiplicative units
$A \oplus B$	plus	$A \& B$	additives
$\mathbf{0}$	zero	$\top$	additive units
$!A$	of course	$?A$	exponentials

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### Focalization

A consequence of this fact is that, **when searching for a proof of  $\vdash \Gamma$ , one can always start by decomposing negative connectives in  $\Gamma$  without losing provability.**

For instance:

- $\vdash \Gamma, (A \wp B) \wp (B \& C)$  is provable
- iff  $\vdash \Gamma, A \wp B, B \& C$  is provable
- iff  $\vdash \Gamma, A \wp B, B$  and  $\vdash \Gamma, A \wp B, C$  are provable
- iff  $\vdash \Gamma, A, B, B$  and  $\vdash \Gamma, A, B, C$  are provable

So without loss of generality, we can assume that any proof of  $\vdash \Gamma, (A \wp B) \wp (B \& C)$  ends like

$$\frac{\frac{\frac{\vdash \Gamma, A, B, B}{\vdash \Gamma, A \wp B, B} (\wp)}{\vdash \Gamma, A \wp B, B \& C} (\&)}{\vdash \Gamma, (A \wp B) \wp (B \& C)} (\wp)$$

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### Any of this two is a focalized proof?

$$\frac{\frac{\frac{\frac{\vdash X, X^\perp}{\vdash X \otimes Y, X^\perp, Y^\perp} ax}{\vdash X \otimes Y, X^\perp \otimes 1, Y^\perp} \otimes}{\vdash (X \otimes Y) \oplus \perp, X^\perp \otimes 1, Y^\perp} \oplus_1}{\vdash \Gamma} \text{cut}$$

$$\frac{\frac{\frac{\frac{\vdash X, X^\perp}{\vdash X, X^\perp \otimes 1} ax}{\vdash X \otimes Y, X^\perp \otimes 1, Y^\perp} \otimes}{\vdash (X \otimes Y) \oplus \perp, X^\perp \otimes 1, Y^\perp} \oplus_1}{\vdash \Gamma} \text{cut}$$

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## Summing-Up Linear Logic rules

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	$\frac{}{\vdash A, A^\perp} \text{ (identity)}$	$\frac{\vdash \Gamma, A \quad \vdash A^\perp, \Delta}{\vdash \Gamma, \Delta} \text{ (cut)}$
<b>Multiplicatives</b>	$\frac{}{\vdash \mathbf{1}} \text{ (one)}$	$\frac{\vdash \Gamma}{\vdash \Gamma, \downarrow} \text{ (false)}$
	$\frac{\vdash \Gamma, A \quad \vdash B, \Delta}{\vdash \Gamma, A \otimes B, \Delta} \text{ (times)}$	$\frac{\vdash \Gamma, A, B}{\vdash \Gamma, A \wp B} \text{ (par)}$
<b>Additives</b>	$\text{ (no rule for zero)}$	$\frac{}{\vdash \Gamma, \top} \text{ (true)}$
	$\frac{\vdash \Gamma, A}{\vdash \Gamma, A \oplus B} \text{ (left plus)}$	$\frac{\vdash \Gamma, A \quad \vdash \Gamma, B}{\vdash \Gamma, A \& B} \text{ (with)}$
	$\frac{\vdash \Gamma, B}{\vdash \Gamma, A \oplus B} \text{ (right plus)}$	
<b>Exponentials</b>	$\frac{\vdash ?\Gamma, A}{\vdash ?\Gamma, !A} \text{ (of course)}$	$\frac{\vdash \Gamma}{\vdash \Gamma, ?A} \text{ (weakening)}$
	$\frac{\vdash \Gamma, A}{\vdash \Gamma, ?A} \text{ (dereliction)}$	$\frac{\vdash \Gamma, ?A, ?A}{\vdash \Gamma, ?A} \text{ (contraction)}$

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## Linear negation

$$(X^\perp)^\perp = X$$

$$(A \otimes B)^\perp = A^\perp \wp B^\perp$$

$$(A \wp B)^\perp = A^\perp \otimes B^\perp$$

$$(A \& B)^\perp = A^\perp \oplus B^\perp$$

$$(A \oplus B)^\perp = A^\perp \& B^\perp$$